ECE 435 – Network Engineering Lecture 9

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12:30pm, Barrows 125

13 February 2024

Announcements

HW#4 was posted (e-mail, DNS)



Last Time – DNS Questions

- DNS Class (usually IN for internet)
 Mostly reserved, with two obsolete networks chaos, hesiod
- Homework: tried for a packet analysis question, grabbed a bunch, turns out decoding a DNS packet is *really* tricky. Original protocol was relatively straightforward but modern real-life there's a lot going on



HW#2 – Programming Notes

- Watch warnings, though I might be running newer version of gcc
- Don't use string operations on binary files
- If no file specified, index.html If no index.html send a 404 error
- ctime prints own linefeed
- If you report HTTP 1.1, don't close connection after file, there might be more requests and you might get "connection reset"



- Be sure to check for unexpected errors what if huge URL is sent?
- Many crashed if I requested the README file. Have to handle unexpected input from user. (in this case, no file extension)
- Traditionally the biggest problem (if the browser refuses to display) is the wrong Content-length:
 If you send less data than you say you will, it will wait forever for it, or else give a "connection reset" if you close the connection.
- Be sure you read everything the browser is sending



(Either big enough buffer, or repeat in loop reading it all). If you send a response before it is done sending it can confuse things. How can you hold an arbitrary size header? malloc()? Do you want to?

- Be sure to drop the leading / in the file part of a URL
- If you use firefox you'll see it might also request favico.ico? Why? What should you return (assuming the file doesn't exist?) 404.



HW#2 – Why write web server in C?

- A pain to write in C.
- But... what language are most webservers written in?
 Apache=C, nginx =C, lighttpd = C



HW#2 – Questions

- browser
 - Error 404 not found
 - Error 418 RFC 2324 coffee protocol (I'm a teapot)
 - Error 451 Unavailable For Legal Reasons / Ray Bradbury
- http header
 - \circ nginx/1.20.1
 - Isn't actually a website, just redirect to the https site



HW#2 – Something Cool

• I do appreciate the pages you made, even if I didn't comment specifically in the grades.



The Transport Layer

	OSI	TCP/IP
7	Application	Application
6	Presentation	
5	Session	
4	Transport	Transport
3	Network	Internet
2	Data Link	Host-to-network
1	Physical	Host-to-network



The Transport Layer

- Responsible for reliable point-to-point data transport independent of whatever lies beneath.
- Sender: receives data from application, breaks up to "segments", adds port number, goes to net layer
- Receiver: re-assembles "segments", passes data to application listening on port
- Provide process-to-process connectivity, and persegment error control and per-flow reliability, as well as rate control



- Can be more reliable than underlying network
- Most common interface "socket" API from homeworks.



Some Transport Layer Protocols

- TCP (Transmission Control Protocol)
 - connection oriented / stateful / per-flow reliability and rate control
- UDP (User Datagram Protocol)
 - stateless / connectionless
- SCTP (stream control transmission protocol)
 - messages like UDP, reliable like TCP
- QUIC
 - running reliable connection over UDP



The Transport Layer

- Terminology: application = process, data-transfer-unit is a segment, traffic is a flow
- addressing each process needs a unique ID. For internet, this is the "port" number (16-bit)
- Rate control
 - Flow control between source and destination
 - Congestion control between source and network
 None in link layer because only one hop?



Can be done by sender or network

 Real time requirements – things like video and audio need extra info such as timestamp, loss rate, etc. So hard to do with raw TCP/UDP



Unreliable, Connectionless – UDP

- User Datagram Protocol (RFC 768)
- Just an 8-byte header tacked onto the data packet
- No reliability, no rate control, stateless
 If you want these things you have to add at higher layer
- Error control optional
- Why none of those things? All add overhead.
 - Used when want packets to get through quickly.
 - Don't care about re-transmits, better for real-time (VOIP, streaming?)



- Easy to implement, for low-level stuff like bootp/dhcp
- Good for broadcasting
- Provides process-to-process communication and persegment error control
- Can send UDP packets to a destination without having to set up a connection first



UDP Header

2 bytes	2 bytes
Source Port	Destination Port
Packet Length	Checksum

- 16-bits: source port (optional, says where it is coming from in case need to respond, 0 if unused)
- 16-bits: destination port
- 16-bits: length (in bytes, includes the header) min: 8, max: 65,515 (<64k to fit in 64k IP packet)
- 16-bits checksum (optional, 0 if unused, see below)
- data follows



Port Number Review

- 16-bit, so 64k of them
- Can map to any you want, but there are certain well-known ones. Look in /etc/services
 For example. WWW is 80/tcp. DNS is 53/udp
- Most OSes, ports <1024 require root (why?)
- 1024 ... 49151 are registered IANA ports
- 49152 ... 65535 are ephemeral ports, dynamic for use by any service



Uniquely identifying flow

- TCP/UDP on IPv4 represented by 104 bit socket pair 5-tuple
 - Source/destination addr
 - Source/destination port
 - protocol ID (TCP or UDP)
- IPv6 in theory has a specific field for this



UDP checksum

- Find info on this in RFC768 and RFC1071
- If set to zero, ignored
- Receiver drops packet if invalid checksum
 Does not request resend, does not notify sender (yes, really, no error message if dropped)



UDP checksum Algorithm

- Ones-complement sum of all 16-bit words in header and payload
- Padded with 0s to be multiple of 16-bits
- Also added to the checksum is the pseudo-header (Layering Violation)
 - Enables receiver to catch problems (delivered to wrong machine) why could this be a problem?
 - IPV4: a 96-bit pseudo header: source IP (2*16), dest



- IP (2*16), protocol (padded to 16), length
- IPv6: 128-bit src IP, 128-bit dest IP, 32-bit UDP len,
 24-bit 0, 8-bit next/type (17 UDP)
- The ones' complement of checksum is put in checksum field. That way when you checksum a valid packet the result will be 0.



Ones' Complement Refresher

- Positive numbers are same as always (with high bit 0)
- Negative numbers are represented by the inverse (bit flipped) of positive. (no adding 1 (that's twos' complement)
- When adding, if there's a carry, it is wrapped around and added in to the low bit ("end-around carry")
- Subtraction is a bit more complicated
- There are two zeros, 0x0000 and 0xffff



Checksum Corner Cases

- What happens if checksum is 0? Conflict with disable checksum? Entered as 0xffff, which in ones complement is -0
- Checksum considered mandatory on IPv6 because IPv6 header not checksummed
- Why would you ever leave checksum out? Takes time to compute, might care about latency over errors [video?]



UDP example Packet

```
0x0000: 8875 563d 2a80 0030 18ab 1c39 86dd 6002 .uV=*..0...9.....
0x0010: 2618 0031 1140 2610 0048 0100 08da 0230
                                                &..1.@&..H....O
0x0020: 18ff feab 1c39 2001 4860 4860 0000 0000
                                                ....9..H'H'....
0x0030: 0000 0000 8844
UDP starts at 0x36:
                                                ....D.9.5.1...W
                       e239 0035 0031 9c0e 8657
0x0040:
        0120 0001 0000 0000 0001 0377 7777 0465
                                                ....www.e
        7370 6e03 636f 6d00 0001 0001 0000 2910
0x0050:
                                                spn.com....).
0x0060: 0000 0000 0000 00
```



UDP example Packet Decoded

- What is source port? What is destination port? Size?
- How can you tell what high-level protocol it is?
 Can you make an educated guess from the ports?



UDP checksum example (from prev slide)

- 16-bit sum of "virtual header" (two IPv6 addresses, protocol (0x0011) and length of udp packet/header (0x0031)) is 0x2'9f8c
- 16-bit sum of UDP header leaving off checksum is 0xe29f
- 16-bit sum of UDP data is 0x2'e1c0
- Add them get 0x6'63eb
- It's a 16-bit sum, so add 0x6 + 0x63eb = 0x63f1 ones complement is 0x9c0e, which matches the UDP checksum field



UDP and the Operating System

- Server: user binds to UDP socket
- OS sets up queue
- Network stack decodes packet at lower level, notes that it is UDP
- Runs checksum, drops it if invalid
- Finds port, looks to see if any processes listening on that port
- If so, adds to queue
- If not, sends an ICMP "port unreachable" error message



 All UDP messages to that port, no matter who sends them, end up in the same queue.



Writing UDP sockets code

- Use SOCK_DGRAM rather than SOCK_STREAM
- Can skip the listen/accept state, as no connection is there. Just receive the packets as they come in.
- Can't read then write, as no connection. For the server to write back to the client it needs to use recvfrom() which also provides ip/port
- To send a packet use sendto()



UDP Socket – Client code



UDP Socket – Server code



Common UDP Services

- Obsolete: echo/discard/users/daytime/quote/chargen
- Nameserver
- bootp/tftp
- ntp (network time protocol)
- Old versions of NFS
- snmp



UDP real-time

- Real-Time Protocol (RFC1889)
- On top of UDP, multiplexes
- data streams
- timestamps

