

ECE 574 – Cluster Computing

Lecture 12

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Announcements

- No Homework over Break
- Did post handout for the project



HW#6 Review

- Coarse

	real	load	convolve	combine	store
1	0m2.013s	173,486	1,535,712	155,207	68,417
2	0m1.321s	144,335	854,183	200,480	69,430

- Fine



fine:	real	load	conv	comb	store
1	0m1.899s	144,764	1,560,142	109,010	39,953
2	0m1.067s	144,665	794,610	56,179	39,717
4	0m0.675s	146,151	364,167	40,649	53,259
8	0m0.486s	145,470	228,161	18,255	43,320
16	0m0.470s	146,099	201,600	13,735	54,268
32	0m0.530s	144,969	284,562	8,026	42,780
64	0m8.107s	144,948	7,833,613	21,380	55,849

- Static/Dynamic with 16 threads

	real	load	conv	comb	store
static	0m0.470s	146,099	201,600	13,735	54,268
dynamic	0m3.776s	146,104	2,318,621	1,189,034	70,071



Shared Memory vs Distributed Systems

Reminder: shared memory has one copy of OS and all programs see one unified memory space.



Shared Memory

- OpenMP is nice to use. But what if your problem won't fit on a single machine?
- How big can a shared-memory machine be?
- SGI UV systems at least 4096 cores and 16TB running one Linux image

<http://www.techeye.net/hardware-2/sgi-builds-pittsburgh-4096-processor-core-16tb-shared-memory-supercomputer>



- Digression about SGI
- Use special NUMA-Linux architecture to spread cache coherence across multiple machines.
- Origin TM and Onyx2 TM Theory of Operations Manual



Linux limitations

- Linux currently maxes out to 4096 or so.
- Somewhat dated “Scaling Linux to the Extreme” paper
problems: cache contention could bring machine to halt
(if a global idle counter, each thread trying to increment
once a second)
lock contention, page cache
- What are the challenges? Locking contention?
- Benefits?



(Relatively) easy to code?

Easier to port code

Many libraries do it for you. For example, OpenBLAS.



Eventually you hit the limit

What's the alternative?



Distributed System

- Communicate over a network
- Many systems each with own memory, communicate via Message passing
- Each node has own copy of operating system
- How do they communicate?



Network Topology

- Packet-switching vs bus
- Ring, mesh, star, line, tree, fully connected
- Cube, hypercube
- Mesh networks and routing
- Routing. Fully connected? Crossbar?



Network Types

- Latency vs Bandwidth
- Top500 in Jun 2015:

interconnect	#
infiniband FDR	160
10GB ethernet	83
infiniband QDR	73
gigabit ethernet	63
Cray Gemini	15

- Ethernet – 10/100/1Gb/10GB/40Gb/s



- InfiniBand – low latency, most common in supercomputers
copper or fiber, GB/s

	SDR	DDR	QDR	FDR-10	FDR	EDR
4x link	8	16	32	40	54	96
12x link	23	48	96	120	163	290

- Cray Gemini – Mesh/torus – 64Gb/s
- Fibrechannel
- Older: custom, Myrinet



Programming a distributed System

- Can you implement by hand?
- Sort of how you can use pthread directly?
- Yes, use ssh (like rsh) to run copy of your program on all machines
- Then write custom network code to open sockets and communicate among them all
- Network code is a pain
- Just crying out for abstraction



Message Passing Interface (MPI)

Abstraction for sending chunks of data around network. You can put together an array of 100 floats, and say "send this to process Y" and like magic it appears there.



MPI

- Message Passing Interface
- Distributed Systems
- MPI 1.0 – 1994. MPI 3.0 – 2012
- MPI 1.2 widely used. MPI2.0 is complicated and adoption not as high as it could be.
- MPICH – CH stands for Chameleon – Argonne and Mississippi State



- MVAPICH – from Ohio State, based on MPICH
- OpenMPI – merger of 3 MPI implementations: FT-MPI from the University of Tennessee, LA-MPI from Los Alamos National Laboratory, and LAM/MPI from Indiana University
- Any other options? PVM was a predecessor
- Python Bindings, Java bindings, Matlab



MPI

Some references

<https://computing.llnl.gov/tutorials/mpi/>

<http://moss.csc.ncsu.edu/~mueller/cluster/mpi.guide.pdf>



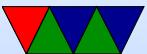
Writing MPI code

- `#include "mpi.h"`
- Over 430 routines
- use `mpicc` to compile
gcc or other compiler underneath, just sets up includes and libraries for you.
- `mpirun -n 4 ./test_mpi`
- `MPI_Init()` called before anything else
- `MPI_Finalize()` at the end
- Error handling – most errors just abort



Communicators

- You can specify communicator groups, and only send messages to specific groups.
- `MPI_COMM_WORLD` is the default, means all processes.



Rank

- Rank is the process number.
- `MPI_Comm_rank(MPI_Comm comm, int size)`
`MPI_Comm_rank(MPI_COMM_WORLD, &rank);`
- You can find the number of processes:
`MPI_Comm_size(MPI_Comm comm, int size)`



Error Handling

- `MPI_SUCCESS` (0) is good
- By default it aborts if any sort of error
- Can override this



Timing

- `MPI_Wtime()`; wallclock time in double floating point.
For PAPI-like measurements
- `MPI_Wtick()`;



Point to Point Operations

- Buffering – what happens if we do a send but receiving side not ready?
- Blocking – blocking calls returns after it is safe to modify your send buffer. Not necessarily mean it has been sent, may just have been buffered to send. Blocking receive means only returns when all data received
- Non-blocking – return immediately. Not safe to change buffers until you know it is finished. Wait routines for



this.

- Order – messages will not overtake each other. Send #1 and #2 to same receive, #1 will be received first
- Fairness – no guarantee of fairness. Process 1 and 2 both send to same receive on 3. No guarantee which one is received



MPI_Send, MPI_Recv

- block – `MPI_Send(buffer, count, type, dest, tag, comm)`
- non-block – `MPI_Isend(buffer, count, type, dest, tag, comm, request)`
- block – `MPI_Recv(buffer, count, type, source, tag, comm, status)`
- non-block – `MPI_Irecv(buffer, count, type, source, tag, comm, request)`
- `buffer` – pointer to the data buffer
- `count` – number of items to send



- type – MPI predefines a bunch. MPI_CHAR, MPI_INT, MPI_LONG, MPI_DOUBLE, etc.
can also create own complex data types
- destination – rank to send it to
- source – rank to receive from. Also can be MPI_ANY_SOURCE
- Tag – arbitrary integer uniquely identifying message.
Can pick yourself. 0-32767 guaranteed, can be higher.
- Communicator – can specify subgroups. Usually use



MPI_COMM_WORLD

- status – status of message, a struct in C
- request – on non-blocking this is a handle to the request that can be queried later to see that status



Fancier blocking send/receives

- Lots, with various type of blocking and buffer attaching and synchronous/asynchronous



Sample code

```
/* MPI Send Example */
#include <stdio.h>
#include "mpi.h"

#define ARRAYSIZE 1024*1024

int main(int argc, char **argv) {

    int numtasks, rank;
    int result, i;
    int A[ARRAYSIZE];
    MPI_Status Stat;
    int count;

    result = MPI_Init(&argc, &argv);
    if (result != MPI_SUCCESS) {
        printf ("Error starting MPI program!.\n");
        MPI_Abort(MPI_COMM_WORLD, result);
    }

    MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
```



```

MPI_Comm_rank(MPI_COMM_WORLD, &rank);

printf("Number of tasks= %d My rank= %d\n",
       numtasks, rank);

if (rank==0) {
    /* Initialize Array */
    printf("Initializing array\n");
    for(i=0; i<ARRAYSIZE; i++) {
        A[i]=1;
    }

    for(i=1; i<numtasks; i++) {
        printf("Sending %d ints to %d\n",
              ARRAYSIZE, i);
        result = MPI_Send(A, /* buffer */
                          ARRAYSIZE, /* count */
                          MPI_INT, /* type */
                          i, /* destination */
                          13, /* tag */
                          MPI_COMM_WORLD);
    }
}
else {

```




```

    result = MPI_Recv(A, /* buffer */
                     ARRAYSIZE, /* count */
                     MPI_INT, /* type */
                     0, /* source */
                     13, /* tag */
                     MPI_COMM_WORLD,
                     &Stat);
    result = MPI_Get_count(&Stat, MPI_INT, &count);
    printf("\tTask %d: Received %d ints from task %d with tag %d \n",
           rank, count, Stat.MPI_SOURCE, Stat.MPI_TAG);
}

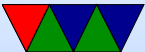
int sum=0, remote_sum=0;

for(i=rank*(ARRAYSIZE/numtasks); i<(rank+1)*(ARRAYSIZE/numtasks); i++) {
    sum+=A[i];
}

if (rank==0) {

    for(i=1; i<numtasks; i++) {
        result = MPI_Recv(&remote_sum, /* buffer */
                          1, /* count */
                          MPI_INT, /* type */

```



```

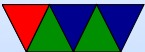
        MPI_ANY_SOURCE, /* source */
        13,             /* tag */
        MPI_COMM_WORLD,
        &Stat);
result = MPI_Get_count(&Stat, MPI_INT, &count);
printf("\tTask %d: (%d) Received %d int from task %d with tag %d \n",
        rank, remote_sum, count, Stat.MPI_SOURCE, Stat.MPI_TAG);
sum+=remote_sum;

}
printf("Total: %d\n", sum);

}
else {
printf("\tRank %d Sending %d\n", rank, sum);
result = MPI_Send(&sum, /* buffer */
        1, /* count */
        MPI_INT, /* type */
        0, /* destination */
        13, /* tag */
        MPI_COMM_WORLD);

}
MPI_Finalize();

```



}

