

# ECE 435 – Network Engineering

## Lecture 15

Vince Weaver

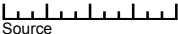
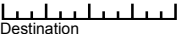
<http://web.eece.maine.edu/~vweaver>

[vincent.weaver@maine.edu](mailto:vincent.weaver@maine.edu)

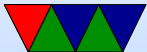
24 October 2017

# Announcements

- Project
- RFC791 post-it note:

Internet Protocol Datagram		<b>RFC791</b>
Source 	Destination 	Version <input type="checkbox"/> <small>If other than version 4, attach form RFC 2460.</small>
<b>Type of Service</b> <input type="checkbox"/> high reliability <input type="checkbox"/> high throughput <input type="checkbox"/> low delay	<b>Precedence</b> <input type="checkbox"/> Routine <input type="checkbox"/> Priority <input type="checkbox"/> Immediate <input type="checkbox"/> Flash <input type="checkbox"/> Flash Override <input type="checkbox"/> CRITIC/ECP <input type="checkbox"/> Internetwork Control <input type="checkbox"/> Network Control	<b>Fragmentation</b> <b>Offset</b> <small>Transport layer use only</small> <input type="checkbox"/> more to follow <input type="checkbox"/> <input type="checkbox"/> do not fragment <input type="checkbox"/> <input type="checkbox"/> this bit intentionally left blank Identifier _____
<b>Protocol</b> <input type="checkbox"/> TCP <input type="checkbox"/> UDP <input type="checkbox"/> Other _____	<b>Length</b> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/> <b>Header Length</b> <input type="checkbox"/> <input type="checkbox"/>	<b>Data</b> <small>Print legibly and press hard. You are making up to 255 copies.</small> _____ _____ _____ _____ _____ _____
<b>Time to Live</b> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/>	<b>Options</b> <small>Do not write in this space.</small>	
<b>Header Checksum</b> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/>		

for more info, check IPv4 specifications at <http://www.ietf.org/rfc/rfc0791.txt>



# HW#6 Review

- Decoding header, much better than last homework
- Network connections
  - CLOSE-WAIT: received a FIN and ACKed it, waiting to close  
Only a few, https and imap
  - ESTAB: established, a few ssh, https, imap connections
  - SYN-RECV: way too many, SYN flood
  - TIME-WAIT: connection closed, waiting a bit before



re-using port

- UNCONN – UDP listening. 789? ipp, mdns (multi-cast DNS, bonjour, can find names on network w/o running DNS), lsof -i udp:789, rpcbind
- LISTEN – listening. Can see ipp (CUPS printing), netbios/microsoft, apparently have SAMBA running,
- Synflood, by default Linux uses SYN cookies to defend against this



# Project

- Can work in groups
- Do something interesting network related.
- Can use any operating system and written in any language (asm, C, python, C++, Java, etc.)
- Coding, benchmarking
- Past projects: network games, firewall config, network attached storage, mesh networks
- Will be a final writeup, and then a 10 minute presentation and demo in front of the class during last week of classes.



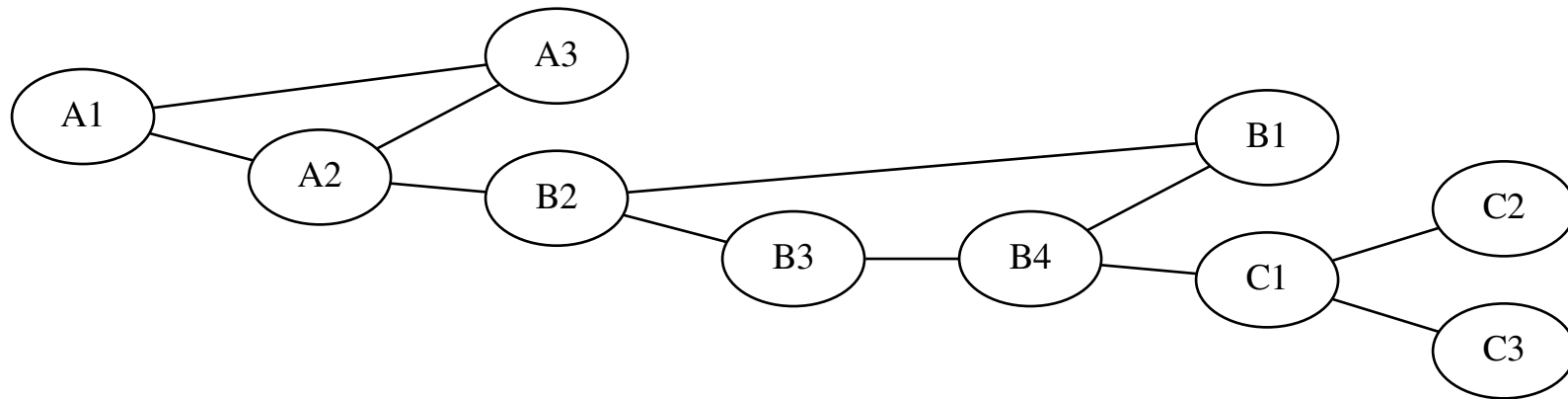
# Hierarchical Routing

- Would you want to have all routers in network on flat network? Crazy routing table
- Split into a hierarchy
- Autonomous System (AS) – a network under control of one group.  
Inside an AS, interior routing, between is exterior routing.
- Systems under same command (same ISP) use intra-domain routing protocol, or interior gateway protocol (IGP)



- Border routers connect to border routers of others
- Inter-domain routing, EGP (exterior gateway protocol)
- types
  - Stub AS – like ISP with customers, one gateway to internet
  - Multihomed AS – multiple gateways (why?)  
redundancy. traffic generally doesn't flow through
  - Transit AS – traffic can flow through network





- Packet A1 - A3 internal A1 - B2 goes to border router and across, then local A1 - C2 goes to border router to B network, across local to B/C border, then finally to C
- If flat network, need to know 10 machines in routing table
- In hierarchical only need to communicate to 2-3 other routers, find way to border router



# Intra-Domain Routing / Interior Gateway Protocols



# RIP (Routing Information Protocol)

- by Xerox, included in BSD, routed RFC 2543
- distance vector routing, with hop count, max 15 hops
  - RIP advertisements over UDP port 52
  - Send advertisement every 30s, or when changes
  - Only sends to neighbors
- Routing table: dest, next hop, distance
- Algorithm
  - Get table update
  - Increment all hops by 1 (you're one hop away)



- Go down list.
  - If route not in table, add it
  - If route there, and next hop same (but cost diff), replace it as this is new info
  - If route there but cost less, replace it
- On power up, comes up with hard-coded routes and values of 1 and no next-hop. Can send packet to request immediate update from neighbors.
- Packet description
- Timers
  - Periodic timer, technically 30s, reality randomized



between 25 and 35 (why?)

- Expiration timer – 180s. If no update in this time, problem, hop count set to 16 (unreachable)
- Garbage collection – 120s – once unreachable, advertise it as such for a while before removing so others notice
- Issues
  - Slow Convergence – a change in routing tables takes 30s per hop to propagate through  
Part of why limited to 15 hops
  - Instability – packets can be caught in loops. Ways to



fix:

Triggered update – send update info immediately, not wait 30s

Split Horizon – if a router sends you update info, don't send this back to it  
Poison Reverse – like split horizon, but when send back, mark as 16 the routes received from that interface.

- There was a RIP2



# OSPF (Open Shortest Path First)

- successor to RIP. RFC2328 (5340 for IPv6)
- Idea of Areas inside of an AS. Split up into areas. Each area connected by backbone router
- Link-state Routing
  - State is flooded: when a change happens (and only then) it sends this state to all neighbors, which send to all neighbors, until the whole network receives it
  - each router uses Dijkstra to find least cost for self, builds table



- Types of link
  - Point-to-point – routers directly connected
  - Transient Link – network with several routers can be simplified?
  - Stub Link – a network connected to only one router
  - Virtual Link – a path between two routers that traverses other routers
- load balancing – supports equal-cost multipath routing (can equally use equal cost routes)
- supports CIDR routing
- support available for multicast



- 8-byte password for authentication
- supports hierarchical
- example? complex!



# Inter-Domain Routing

- Can be complicated.
- Say company with network, and two connections to outside X,Y. Don't want to send packets out and back even if it looks like lower cost.
- Also don't want to transit packets between X and Y for outsiders. Policy.



# BGP (border gateway protocol)

- Intro in 1989, four versions – BGP4 RFC 4271
- Uses TCP (reliable) port 179.
- Works for both IPv4 and IPv6 (the latter as an extension)
- Uses path vector rather than distance vector
  - full path, not just next-hop
  - exchanges info with neighbor, but includes complete path info to avoid looping.
  - Each AS has unique number, so if it sees itself in the path knows there is a loop.



- Policy routing – can also reject new route based on policy
- Four types of messages – open, update, keepalive, notification
- Whole table not passed around (Due to size), only updates
- Due to size of internet, uses distance vector over link state.
- interior and exterior BGP
- iBGP makes sure that the setups for multiple gateway routers are kept synchronized



- eBGP used to talk between other exterior routers at peers.
- Keeps track of all feasible paths, but only advertises the “best” one



# Routing table Size

- Example. Full BGP of internet backbone router might have more than 300,000 entries (2010) now over 700,000.
- <http://bgp.potaroo.net/>
- Some routers had limit of 512k so on August 12 2014 part of internet went down when crossed the border.
- Ipv6 currently only around 20k



# Peering

- How companies agree to connect their networks together. There's not really a master connection, but instead companies agree to have routers talk to each other via BGP.
- Transit – pay money to pass through network.
- Peering – In many cases no money changes hands. Why? Well if you have a lot of users, but no content, people won't stay with you. Same if you have content but no



access to users. Averages out and is mutually beneficial.

- Increased redundancy
  - Increased capacity
  - Increased routing control
  - Improved performance
  - Fame (high-tier network)
  - Ease of requesting aid (?)
- Customer – you buy an internet connection
  - Peering locations, often in large data centers.  
At one point there were 4 major ones (Metropolitan Area



Exchange) MAE-East (Virginia) [in basement of parking garage, at one point half of internet went through here], Chicago, NY, SF. All defunct now

- Depeering – if you think you aren't getting a good deal, break up. Some situations there is a fight, a hope that the customers lose enough performance will have to repeer.
- Related – net neutrality
- Tiers – Tier 1 network is one that can reach rest of internet without paying for transit; Tier 2 peers with



some but purchases for other; Tier 3 only purchases



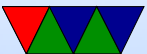
# Routing Security Issues

- Problems – routing black hole, use BGP to send addresses intentionally to 0.0.0.0 and get dropped. BGP will propagate
- router update mistakes can accidentally blackhole parts of the internet
- In 2008 Pakistan was trying to blackhole Youtube and accidentally announced to world via BGP and took it down world wide



# Implementations

- Actual Router
- Can install on your Linux machine
- Zebra was traditionally, discontinued
- Quagga
- BIRD
- OpenBGPD and OpenSPFD



- Potentially dangerous to mess around with unless you isolate your network well

