

ECE 271 – Microcomputer Architecture and Applications Lecture 9

Vince Weaver

`http://web.eece.maine.edu/~vweaver`

`vincent.weaver@maine.edu`

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Announcements

- Read Chapter 3.5 to 3.7
- Lovebyte Update
- Lab schedule update

Finish Lab#3 (keypad) as soon as possible

Start Lab#4 as soon as possible

Next week is catch-up lab to finish any outstanding labs

Note no Monday lab on Monday due to President's day



Lab #4 Notes

- Redo Lab#1, but in Assembly language



Lab #4 Notes

We want to convert the following C to assembly

```
RCC->AHB2ENR |= RCC_AHB2ENR_GPIOBEN;
```

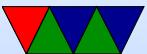
A lot is going on in provided code. Some low-level C and preprocessor directives are used to make RCC be a pointer to address 0x40021000. A cast is used to make this pointer be of type RCC_TypeDef which has the offsets for the various sub-registers (see next slide)

```
#define PERIPH_BASE      ((uint32_t)0x40000000)
#define AHB1PERIPH_BASE (PERIPH_BASE + 0x00020000)
#define RCC_BASE        (AHB1PERIPH_BASE + 0x1000)
#define RCC              ((RCC_TypeDef *) (RCC_BASE))
```



Lab #4 Notes

- Do we remember how structs work in C?
- What's the difference between `a.b` VS `a->b`
- The latter is indexing from a pointer



Lab #4 Notes

The RCC_Typedef is in the provided stm32l476xx.h

```
typedef struct {  
    __IO uint32_t CR;           // Control Register, offset 0x00  
    ....  
    __IO uint32_t AHB2ENR;     // AHB2 periph control, offset 0x4c  
    ....  
} RCC_Typedef;
```

The key understanding is that we want to access the 32-bit value that's at offset 0x4c from the beginning of the base



Lab #4 Notes

The code that implements the register setting can be done like this:

```
ldr    r1,=RCC_BASE           ; r1=&RCC
ldr    r3,[r1,#RCC_AHB2ENR]   ; r3=RCC->AHB2ENR
orr    r3,#RCC_AHB2ENR_GPIOBEN ; r3 = r3 | RCC_AHB2ENR_GPIOBEN
str    r3,[r1,#RCC_AHB2ENR]   ; RCC->AHB2ENR = r3
```

We provide defines in assembly which you can use for value/masks rather than having to do the raw hex codes.



Clearing values

- use the BIC instruction to clear bits

Why not just use AND? Because a mask has lots of 1s might not fit in available constant room

```
and r3,#0xfffffffffe      ; wont fit in instruction, too big
bic r3,#1                  ; same as: and r3,#~1 but fits in constant
```

- You can usually include complicated C-stye constant manipulations, things like

```
and r3,r3,#(GPIOBEN + 0x5 | (1<<3))
```

though note that the Keil assembler might not support the full range



Something-cool Notes / BSRR

- Using the BSRR register to set/reset GPIO pins without having to do a read/modify/write.
- Atomic operation?
- Bit set/reset Register.
 - Write a bit pattern of 0 or 1
 - 0 means leave alone
 - In bottom 16-bits, 1 means set that GPIO to 1
 - In the top 16-bits, 1 means reset that GPIO to 0



Assembler – Code comments

- Can use C and C++ style comments
- Keil: can use ; (makes rest of line a comment)
- Linux/gas: Can use @ for beginning of line



Functions/Subroutines

- Why use them?



Sample C

```
int sum(int a, int b) {  
    return a+b;  
}
```

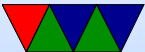
```
int main(int argc, char **argv) {  
    int result;  
    result=sum(1,2);  
}
```



Sample Assembly

```
sum
    add    r0,r1,r2    ; result=arg1+arg2
    bx    lr          ; jump to saved address in link register

main
    mov    r0,#1       ; set arg1 to 1
    mov    r1,#2       ; set arg2 to 2
    bl     sum         ; call sum function, put current
                    ; program counter+4 into link register
```



Subroutines on ARM

- b1 branch and link instruction
 - Sets the link register LR (r14) as the memory address of the next instruction immediately after the BL ($PC+4$ on Thumb-2)
 - Adjust the PC to be the memory address of the first location of where we want to transfer execution
- After executing, LR has the return address



Returning from a Subroutine

- Use the BX LR instruction, which says to branch to the address located in the LR register. (the X is for exchange; historical THUMB reasons)



Saving/Restoring values in functions

- To preserve registers at start of function can

```
push {r0, r1, lr}
```

- You will want to save the link register if not a leaf function (meaning, you call another function from inside)
- At end you can

```
pop {r0, r1, lr}
```

before using `bx lr` to return

- Alternately, if the LR register was pushed on the stack, you can use the clever hack of doing

```
pop {r0, r1, pc}
```



to pop the link register directly into the PC to return without the extra branch

; p164



The ABI – The Application Binary Interface

- A Document, produced often by the processor maker
- An agreement of how functions / code talk to each other
- A common standard so compilers, libraries, and code can call each other and know how to set things up
- Useful to have for your own code. Might be slightly less efficient, but better than for every function you call having to save/setup different registers
- What kinds of things are included?
 - What registers to put things in? Register allocation?

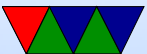


- Alignment of stack (4 bytes? 8 bytes?)
- How to pass 8/16/32/64 byte values
- How to pass floating point values
- Where does the return value go?
- System calls
- Frame pointer



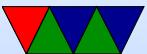
ARM ABI

- On Linux there have been at least 4
- armbe – big endian
- armle – little endian
- EABI – extended (new) ABI
- armhf – EABI but fancier (hard) floating point support



Calling Conventions

- r0/r1/r2/r3
 - parameters/scratch
 - caller saved, so if you want the value in say r3 to be the same after a function call you have to save it/restore it to memory
 - r0/r1 also used as return value from function
- r4/r5/r6/r7/r8/r10/r11 = variables
 - callee saved. You can count on this having the same value after a function as before. If you are in a function



and want to use it, must save/restore it. Often this done at function entry/exit

- r9 – implementation dependent (thread-local register?)
- r12 = reserved by linker?
- r13 = stack pointer
- r14 = LR (link register)
- r15 = PC (program counter)



Calling Conventions – Corner Cases

- Return value in r0. Might be in r1 or more if bigger than 32 bits
- What happens if more than 4 arguments?
- What happens if more than 32-bits (use 2 registers, even/odd for 64-bits)



Calling Conventions – Corner Cases

- How do you pass something complicated like an array or struct?
- Call by value or by reference
- Can pass a pointer in a 32-bit register



Disassembler

```
#include <stdio.h>
int i;
int main(int argc, char **argv) {
    for(i=0;i<100;i++) {
        printf("Hello!\n");
    }
    return 0;
}
```

```
gcc -Wall -mthumb -march=armv7-a -o test test.c
```

Disassembly of section .bss:

```
0002102c <i>:
    2102c:      00000000      andeq    r0, r0, r0
```

Disassembly of section .rodata:

```
000104fc <_IO_stdin_used>:
    104fc:      00020001      andeq    r0, r2, r1
```



```
10500:      6c6c6548      cfstr64vs      mvdx6, [ip], #-288      ; 0xffffffff0
10504:      Address 0x00010504 is out of bounds.
```

0001043c <main>:

```
1043c:      b580          push    {r7, lr}
1043e:      b082          sub     sp, #8
10440:      af00          add     r7, sp, #0
10442:      6078          str     r0, [r7, #4]
10444:      6039          str     r1, [r7, #0]
10446:      f241 032c    movw   r3, #4140      ; 0x102c
1044a:      f2c0 0302    movt   r3, #2
1044e:      2200          movs   r2, #0
10450:      601a          str     r2, [r3, #0]
10452:      e010          b.n    10476 <main+0x3a>
10454:      f240 5000    movw   r0, #1280     ; 0x500
10458:      f2c0 0001    movt   r0, #1
1045c:      f7ff ef42    blx    102e4 <puts@plt>
10460:      f241 032c    movw   r3, #4140     ; 0x102c
10464:      f2c0 0302    movt   r3, #2
10468:      681b          ldr     r3, [r3, #0]
1046a:      1c5a          adds   r2, r3, #1
1046c:      f241 032c    movw   r3, #4140     ; 0x102c
10470:      f2c0 0302    movt   r3, #2
10474:      601a          str     r2, [r3, #0]
```



```
10476:      f241 032c      movw    r3, #4140          ; 0x102c
1047a:      f2c0 0302      movt    r3, #2
1047e:      681b          ldr     r3, [r3, #0]
10480:      2b63          cmp     r3, #99 ; 0x63
10482:      dde7          ble.n   10454 <main+0x18>
10484:      2300          movs   r3, #0
10486:      4618          mov    r0, r3
10488:      3708          adds   r7, #8
1048a:      46bd          mov    sp, r7
1048c:      bd80          pop    {r7, pc}
```



Assembler directives

- AREA – declare a new area (code/data/bss)
AREA myData, DATA, READWRITE
- ENTRY – declare entry point into the code
You might think this is “main” on C, but actually it is usually something called `_start`, a lot of things happen before `main()` gets called
- ALIGN – align the current memory address for performance, or some things (like variables on stack) must be aligned



- DCB – reserve space for bytes
array DCB 1,2,3,4
hello DCB "Hello World!",0
- DCW – reserve space for 16-bit values
- DCD – reserve space for 32-bit values
- DCFS/DCFB – floating point
- SPACE – restore unreserved data (BSS)
p SPACE 255 ; reserve 255 zeros
- FILL – reserve space and fill it with a value
f FILL 20,0xff,1 ; allocate 20 bytes, fill with 0xff
- EQU – like #define in C, lets you set constants.



MAXCPUS EQU 8

On Linux same, but `.equ MAXCPUS, 8`

- RN – alias a register name, if you want to use something like X instead of R3
- EXPORT/IMPORT – export says to make symbol globally visible (`.globl` on Linux). Import says a symbol is external, like “extern” on C.
- INCLUDE/GET – like include directive in C, includes another file when assembling
- PROC/ENDP – start and end of function. Mostly to make debugging easier? Doesn't actually change



generated code?

