

ECE 271 – Microcomputer Architecture and Applications Lecture 14

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Announcements

- Read Chapters 8 + 9
- Midterm, Tuesday, 8 March



Midterm Review

- Closed book/notes.
- Short answer.
- Will be on mostly C and assembly language
- I expect you to know at least basic C
- This includes being able to set/clear bits using the bitwise logic operations
- Assembly language, I will provide a table of THUMB2 instructions so no need to memorize.
- No need to memorize all of the MODER register fields,



etc.

- Things like, what does this code do? Or, add comments to this code, or, what is wrong with this code.
- Know high-level things from lab, GPIO, LCD, Scanning, Stepper
- Basic understanding of what the hardware is doing, but not super detailed
- Questions may be similar to those on pre/post-lab
- There will be a question on setting the N/Z/C/V flags in assembly language. A sample problem of this type was sent out.



- Also stack. Where do local vars go?
- ABI, know why we have one. Where args go.



N/Z/C/V review

- We went over some examples from the sample problem handout
- Remember, ARM only updates the flags if "S" is specified in instruction (e.g. `adds`). Also `cmp` always sets flags even if no S
- N (negative) flag is set if result is negative (most-significant bit (MSB) is 1)
- Z (zero) flag is set if result is zero
- C (carry) flag is set if carry out of MSB when



addition/subtraction. Also gets set to shifted out bit in shift and rotate instructions

- V (overflow) flag is set if carry into the MSB is not the same as the carry out of the MSB



Lab #6 Information

- Re-do the stepper motor lab (lab#5), but in assembly
- Mostly learning to write functions in assembly



Lab #6 – The Delay Code (C)

```
for (i=0; i<6000; i++) ;
```



Lab #6 – The Delay Code (asm)

A quick straightforward translation

```
    mov     r5,#0           // load 0 into iterator
delay_loop
    // loop body usually goes here, which is empty
    add     r5,r5,#1       // increment iterator
    cmp     r5,#6000       // compare to 6000
    bne     delay_loop     // if we aren't there yet, loop
```



Lab #6 – The Delay Code (optimized)

A somewhat shorter version. Note that optimizing delay code can be ill-advised because making it faster changes the timing.

```
    mov     r5,#6000      // change to count down to 0
delay_loop
    subs   r5,r5,#1      // subtract and update flags
    bne   delay_loop    // branch if Zero flag not set
```



Lab #6 – The Delay Code (gcc Compiler)

Note this is with no optimizations, so it has to save out `i` to memory

```
    mov    r3,#0           // load 0 into i
    str    r3,[fp,#-8]     // i is local var on stack
                                // indexed off of frame pointer
    b     next           // skip ahead
loop
    ldr    r3,[fp,#-8]     // load i from local var area
    add    r3,r3,#1       // increment i
    str    r3,[fp,#-8]     // store i back to local var
next
    ldr    r2,=5999       // load 5999
    cmp    r3,r2          // see if we've reached it
    ble   loop           // if not, loop
```



Lab #6 – The Delay Code (gcc Compiler, Optimized)

- With -O2 optimization the compiler realized that the loop didn't do anything and just skipped it all together
- How can you avoid this happening?
 - Declare i to be volatile. This works, but it forces the unnecessary loads/stores to memory
 - Could try putting an inline assembly call to a nop



Lab #6 – Converting to a function

- Add a label for the entry point
- Grab the arguments from r0...r3
- Save any callee saved registers you use (r4...r12) on stack
- Save link register (lr) if you're not a leaf function
- Advanced: allocate any local vars if needed
- At end, restore values off stack
- Put return value in r0 if there is one
- Return from the function



Lab #6 – Delay converted to Function

```
// Delay, with amount in r0
delay PROC // entry point
    push    {r5} // save r5 (callee saved register)
            // we are a leaf function so
            // we don't need to save lr
    mov     r5,#0 // set loop iterator to 0
delay_loop
    add     r5,r5,#1 // increment iterator
    cmp     r5,r0 // compare to the passed in argument
    bne     delay_loop // if we aren't there yet, loop
    pop     {r5} // restore r5 from stack
    bx     lr // return to callsite
ENDP // marks end of procedure (for Keil)
```



Lab #6 – Calling a function (asm)

- If for whatever reason you need caller-saved registers (r0..r3) to be preserved, save them on the stack first
- Next, set up the arguments
- Then branch and link to the function

```
mov    r0, #6000    // load 6000 as first argument
bl     delay        // branch-and-link to delay
                        // stores value of next instruction
                        // link-register so we can return
                        // to right place
```



lab #6 – Using Arrays in Assembly

Code for doing a step sequence might look like this

```
int steps[4]={0x00480084,0x00880044,0x00840048,0x00440088};
int current_step,i;

for(i=0;i<4;i++) {
    current_step=steps[i];
    GPIOB->BSRR=current_step;
    delay(6000);
}
```



lab #6 – Using Arrays in Assembly

```
mov     r5 ,=GPIOB           // point to GPIOB region
mov     r6 ,#0               // setup loop iterator

loop
ldr     r1 ,=steps           // point to steps array
ldr     r2 ,[r1,r6, LSL 2]   // index into steps array with iterator
//     LSL 2 multiplies by 4 as we're an int
//     (4-byte) value

str     r2 ,[r5 ,#BSRR]     // store out to GPIOB->BSRR
mov     r0 ,#6000           // load value to delay
bl      delay               // call delay function
add     r6 ,r6 ,#1          // increment iterator
cmp     r6 ,#4              // see if done
bne     loop                // if not, loop

steps                                     // define an array of four 32-bit values
DCD    0x00480084 ,0x00880044 ,0x00840048 ,0x00440088
```

