# ECE 571 – Advanced Microprocessor-Based Design Lecture 2

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# Announcements

- I've added some optional readings to the website if you want to review Computer Architecture a bit. You can access the 2007 Edition of Patterson and Hennessy for free via the UMaine Library website.
- HW#1 will be posted
- I am handing out account username/passwords for the homework.
- Accounts: Log in to the haswell "haswell-ep" machine for homework. Make sure you connect to port 2131.



ece571-1 names are a bit impersonal.

Use passwd to change your password.

You can use chfn to change your name as it appears in w if you want.

Please use the accounts wisely



# **Evaluating Performance of Modern Systems**



# **Benchmarks**

- When measuring performance, need a reference workload to compare
- Ideally reproducible, portable, easy to compile, relevant
- Benchmarks can be gamed



# Selected Commonly Seen Benchmarks

- SPEC
  - CPU 2000, CPU 2006, CPU 2017 Commercial, Single-threaded (floating point and integer)
  - OMP Commercial, Parallel
  - ∘ jbb Java
- HPC Challenge Free. HPL (Linpack). Highperformance / Linear Algebra
- HPCG (conjugate gradient) new replacement for HPL
- PARSEC Free, Multithreaded / CMP



- MiBench Free, Embedded (2000)
- BioBench, BioParallel Free, Bio/Data-Mining
- Imbench Free, Operating System
- iobench Disk I/O
- Stream Memory



### **Measuring Performance**



# Using time

#### • For example

<pre>\$ time xhpl</pre>						
• • •						
real	0m9.484s					
user	0m29.150s					
sys	0m7.440s					

### Real Time = Wall clock

User Time = Time used by program alone

Sys Time = Time used by OS

• When could Real be greater than User?



Other users/jobs on system. When could User be greater than Real? Multiple threads.

• Run multiple times and notice time changes



## What if Time isn't Enough?



# What are Hardware Performance Counters?

- Registers on CPU that measure low-level system performance
- Available on most modern CPUs; increasingly found on GPUs, network devices, etc.
- Low overhead to read



# Hardware Implementation of Counters

- Not much documentation available
- Jim Callister/Intel: "Confessions of a Performance Monitor Hardware Designer" 2005, Workshop on Hardware Performance Monitor Design
  - Transistors free, wires not. Also design time, validation, documentation, time to market. PMU has tentacles "everywhere" bringing data back to center.
  - Architect too much, lower performance, events don't



map well to hardware. Architect too little.. software design harder.

- Which events are important? Are cache misses important if don't hurt performance? (no stalls)
- Mapping events to signal difficult. On critical path.
   Not enough wires. Combining signals hard if distance between wires.
- Use logging. May miss events in "shadow" of another event being logged. Use random behavior?



# **Learning About the Counters**

- Number of counters varies from machine to machine
- Available events different for every vendor and every generation
- Available documentation not very complete (Intel Vol3b, AMD BKDG, ARM ARM/TRM)



### Low-level interface

- on x86: MSRs
- ARM: CP15 system control register



# Overflow

- overflows after hitting a threshold (often when wrapping, most counters are between 32 and 44 bits wide)
- One use is to keep track of counters that may wrap multiple times between reads
- If want to overflow earlier, init to a high value. So 0xc0000000 to overflow at 1 billion



# Accuracy, Determinism vs Overcount

- Determinism same count every time you run
- Overcount an event counts more than the expected amount



# **SW Sources of Non-Determinism**

- Accessing changing values, such as time
- Pointer-value dependencies



# Linux interface

- Abstract away.
- perf\_event\_open(). See the manpage.
- Very complicated system call.
- Most people use perf or PAPI rather than calling it directly.



# perf tool

A a tutorial on perf can be found here: https://perf.wiki.kernel.org/index.php/Tutorial



# perf list

Lists available events List of pre-defined events (to be used in -e): cpu-cycles OR cycles instructions cache-references cache-misses branch-instructions OR branches branch-misses bus-cycles

cpu-clock
task-clock
page-faults OR faults
minor-faults
major-faults
context-switches OR cs

[Hardware event] [Hardware event] [Hardware event] [Hardware event] [Hardware event] [Hardware event] [Hardware event]

[Software event] [Software event] [Software event] [Software event] [Software event] [Software event]



## perf stat – Aggregate results

vince@arm:~/class/ece571\$ perf stat ./matrix\_multiply
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix\_multiply':

11585.144036	task-clock	#	0
19	context-switches	#	0
0	CPU-migrations	#	0
1,633	page-faults	#	0
10,343,746,076	cycles	#	0
5,031,717	stalled-cycles-frontend	#	0
9,521,135,479	stalled-cycles-backend	#	92
1,176,286,814	instructions	#	0
		#	8
137,835,961	branches	#	11
831,736	branch-misses	#	0

- 0.999 CPUs utilized
- 0.000 M/sec
- # 0.000 M/sec
- # 0.000 M/sec
- # 0.893 GHz
- 0.05% frontend cycles idle
- 92.05% backend cycles idle
- 0.11 insns per cycle
- # 8.09 stalled cycles per insn
- # 11.898 M/sec
- # 0.60% of all branches

11.591796875 seconds time elapsed



# perf stat – Specifying Events

vince@arm:~/class/ece571\$ perf stat -e instructions,cycles ./matrix\_multip
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix\_multiply':

1,174,788,622	instructions	#	0.14	insns	per	cycle
8,346,588,065	cycles	#	0.000	GHz		

12.394775391 seconds time elapsed



# perf stat – Specifying Masks

:u is user, :k kernel ARM Cortex A9 cannot specify this distinction (results shown here are x86)

vince@arm:~/class/ece571\$ perf stat -e instructions,instructions:u ./matri
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix\_multiply':

950,526,051 instructions#0.00 insns per cycle945,661,967 instructions:u#0.00 insns per cycle

1.052072277 seconds time elapsed



### <u>libpfm4 – Finding All Event Names</u>

./showevtinfo Supported PMU models: [51, perf, "perf\_events generic PMU"] [65, arm\_ac8, "ARM Cortex A8"] [66, arm\_ac9, "ARM Cortex A9"] [75, arm\_ac15, "ARM Cortex A15"] Detected PMU models: [51, perf, "perf\_events generic PMU", 80 events, 1 max encoding, 0 counters, OS g [66, arm\_ac9, "ARM Cortex A9", 57 events, 1 max encoding, 2 counters, core PMU] Total events: 254 available, 137 supported . . . #-----IDX : 138412068 PMU name : arm\_ac9 (ARM Cortex A9) Name : NEON\_EXECUTED\_INST Equiv : None Flags : None Desc : NEON instructions going through register renaming stage (approximate) Code : 0x74#------. . . .



# libpfm4 – Finding Raw Event Values

```
./check_events NEON_EXECUTED_INST
Supported PMU models:
[51, perf, "perf_events generic PMU"]
[65, arm_ac8, "ARM Cortex A8"]
[66, arm_ac9, "ARM Cortex A9"]
[75, arm_ac15, "ARM Cortex A15"]
Detected PMU models:
[51, perf, "perf_events generic PMU"]
[66, arm_ac9, "ARM Cortex A9"]
Total events: 254 available, 137 supported
Requested Event: NEON_EXECUTED_INST
Actual Event: arm_ac9::NEON_EXECUTED_INST
PMU : ARM Cortex A9
<u>IDX</u> : 138412068
Codes
        : 0x74
```



# perf – Using Raw Event Values

vince@arm:~/class/ece571\$ perf stat -e r74 ./matrix\_multiply
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix\_multiply':

1 r74

11.303955078 seconds time elapsed



# perf stat – multiplexing

perf stat -e instructions, instructions, branches, cycles, cycles ./matrix\_multiply Matrix multiply sum: s=27665734022509.746094 Performance counter stats for './matrix\_multiply': 1,178,121,057 instructions # 0.12 insns per cycle [40.23%] 0.12 insns per cycle [60.25%] 1,180,460,368 instructions # 138,550,072 branches [80.09%] 0.000 GHz 9,999,614,616 cycles # [79.85%] 9,926,949,659 cycles # 0.000 GHz [20.17%] 11.214630127 seconds time elapsed

Note same event not same results, approximate because an estimate. Percentage shown is percentage event was active during run.



### perf stat – all cores

vince@arm:~/class/ece571\$ sudo perf stat -a ./matrix\_multiply
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix\_multiply':

24089.660644	task-clock	#	2.001 CPUs utilized	[100.00%]
105	context-switches	#	0.000 M/sec	[100.00%]
1,641	page-faults	#	0.000 M/sec	
9,218,451,619	cycles	#	0.383 GHz	[100.00%]
9,707,195	stalled-cycles-frontend	#	0.11% frontend cycles idle	[100.00%]
8,393,095,067	stalled-cycles-backend	#	91.05% backend cycles idle	[100.00%]
1,193,164,945	instructions	#	0.13 insns per cycle	
		#	7.03 stalled cycles per insn	[100.00%]
139,913,572	branches	#	5.808 M/sec	[100.00%]
1,221,237	branch-misses	#	0.87% of all branches	

12.040527344 seconds time elapsed

Run on *all* cores of system even if your process not running there. –a option. Need root permissions. (Why? Security)



# perf record – sampling

```
vince@arm: ~/class/ece571$ time ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
real0m10.747s
user0m10.688s
sys0m0.055s
vince@arm: ~/class/ece571$ time perf record ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
[ perf record: Woken up 2 times to write data ]
[ perf record: Captured and wrote 0.454 MB perf.data (~19853 samples) ]
real0m12.009s
user0m11.797s
sys0m0.203s
```

perf record creates perf.data, use -o to specify output



# perf report – summary of recorded data

99.62%	matrix_multiply	matrix_multiply	[.] naive_matrix_multiply
0.38%	<pre>matrix_multiply</pre>	[kernel.kallsyms].head.text	[k] 0xc0046a54
0.00%	matrix_multiply	ld-2.13.so	[.] _dl_relocate_object
0.00%	matrix_multiply	[kernel.kallsyms]	[k]do_softirq

Our benchmark is simple (only one function) so the profiled results are not that exciting.

The [k] indicates that profile happened while the kernel was running.



# Similar ways to get Similar Results

- Valgrind/Callgrind valgrind - -tool=callgrind BENCHMARK then run callgrind\_annotate Note Valgrind is probably around 50 times slower
- Use gprof

Compile your code with -pg Run *gprof BENCHMARK* 



# perf annotate – show hotspots in assembly

0.00	:	845a:	vldr	d7,	[pc, #124]	;	84d8 <naive_matrix_m< td=""></naive_matrix_m<>
30.97	:	845e:	adds	r1, 1	r4, r3		
1.43	:	8460:	add.w	r3, 1	r3, #4096	;	0x1000
1.17	:	8464:	adds	r2, a	#8		
1.36	:	8466:	cmp.w	r3, a	#2097152	;	0x200000
2.97	:	846a:	vldr	d5,	[r2]		
2.62	:	846e:	vldr	d6,	[r1]		
2.78	:	8472:	mov	r9, 1	r2		
2.42	:	8474:	vmla.f64	4	d7, d5,	d6	3
53.81	:	8478:	bne.n	845e	<naive_matr< td=""><td>i&gt;</td><td>c_multiply+0x72&gt;</td></naive_matr<>	i>	c_multiply+0x72>
0.01	:	847a:	adds	r5, a	#1		

The annotated results show a branch and an add instruction accounting for 83% of profiles. Likely this is due to skid and the key instruction is the previous vmla.f64 floating point multiply instruction. The processor just isn't able to stop at the exact instruction when the interrupt comes in.



# Skid

- Beware of "skid" in sampled results
- This is what happens when a complex processor cannot stop immediately, so the reported instruction might be off by a few instructions.
- Some processors do not have this problem, other Intel processors have special events that can compensate for this.

