

# **ECE 571 – Advanced Microprocessor-Based Design Lecture 24**

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# Project/HW Reminder

- Lecture delayed due to snow
- Homework #11 was posted/due



# HW#7, TLB

- Why not more pagefaults? Kernel stat wrong? Prefaulting by kernel?

Odd major faults always 0

```
/usr/bin/time -v
```

major only means goes to disk, in our case we are allocating memory and filling it so no going to disk

So why not paging of executable in? Well the first time you run it after boot there should be, but after that likely in disk cache so not need to go to disk.



- Note a TLB miss does not always equal a page fault
- Number of pages to cover 24MB, true that in theory also needs some 2nd-level pages too (hard to know how many without knowing where in memory things are)



# Graphics and Video Cards



# Old CRT Days

- Electron gun
- Horizontal Blank, Vertical Blank
- Atari 2600 – only enough RAM to do one scanline at a time
- Apple II – video on alternate cycles, refresh RAM for free
- Bandwidth key issue. SNES / NES, tiles. Double buffering vs only updating during refresh



# Old 2D Video Cards

- Framebuffer (possibly multi-plane), Palette
- Dual-ported RAM, RAMDAC (Digital-Analog Converter)
- Interface (on PC) various io ports and a 64kB RAM window
- Mode 13h
- Acceleration – often commands for drawing lines, rectangles, blitting sprites, mouse cursors, video overlay



# Modern Graphics Cards

- Can draw a lot of power
- 2D (optional these days)
- 3D
- Video decoders





# Interface

- Integrated or stand alone
- Integrated traditionally less capable, but changing. Share Memory bandwidth, take memory.



# GPUs

- Display memory often broken up into tiles (improves cache locality)
- Massively parallel matrix-processing CPUs that write to the frame buffer (or can be used for calculation)
- Texture control, 3d state, vectors
- Front-buffer (written out), Back Buffer (being rendered)  
Z-buffer (depth)
- Originally just did lighting and triangle calculations. Now shader languages and fully generic processing



# Video RAM

- VRAM – dual ported. Could read out full 1024Bit line and latch for drawing, previously most would be discarded (cache line read)
- GDDR3/4/5 – traditional one-port RAM. More overhead, but things are fast enough these days it is worth it.
- Confusing naming, GDDR3 is equivalent of DDR2 but with some speed optimization and lower voltage (so higher frequency)



# Busses

- DDC – i2c bus connection to monitor, giving screen size, timing info, etc.
- PCIe (PCI-Express) – most common bus in x86 systems  
Original PCI and PCI-X was 32/64-bit parallel bus  
PCIe is a serial bus, sends packets  
Can power 25W, additional power connectors to supply  
can have 75W, 150W and more  
Can transfer 8GT/s (giga-transfers) a second  
In general PCIe is limiting factor to getting data to GPU.



# Connectors

CRTC (CRT Controller) Can point to same part of memory (mirror) or different.

- RCA – composite/analog TV
- VGA – 15 pin, analog
- DVI – digital and/or analog. DVI-D, DVD-I, DVD-A
- HDMI – compatible with DVI (though content restrictions). Also audio. HDMI 1.0 – 165MHz, 1080p



or 1920x1200 at 60Hz. TMDS differential signaling. Packets. Audio sent during blanking.

- Display Port – similar but not the same as HDMI
- Thunderbolt – combines PCIe and DisplayPort. Intel/Apple. Originally optical, but also Copper. Can send 10W of power.
- LVDS – Low Voltage Differential Signaling – used to connect laptop LCD



# LCD Displays

- Crystals twist in presence of electric field
- Asymmetric on/off times
- Passive (crossing wires) vs Active (Transistor at each pixel)
- Passive have to be refreshed constantly
- Use only 10% of power of equivalent CRT



- Circuitry inside to scale image and other post-processing
- Need to be refreshed periodically to keep their image
- New “bistable” display under development, requires no power to hold state





# Interfaces

- OpenGL – SGI
- DirectX – Microsoft
- For consumer grade, driven by gaming



# GPGPUs

- Interfaces needed, as GPU companies do not like to reveal what their chips do at the assembly level.
  - CUDA (Nvidia)
  - OpenCL (Everyone else) – can in theory take parallel code and map to CPU, GPU, FPGA, DSP, etc
  - OpenACC?



# Other Accelerator Options

- XeonPhi – came out of the larabee design (effort to do a GPU powered by x86 chips). Large array of x86 chips(p5 class on older models, atom on newer) on PCIe card. Sort of like a plug-in mini cluster. Runs Linux, can ssh into the boards over PCIe. Benefit: can use existing x86 programming tools and knowledge.
- FPGA – can have FPGA accelerator. Only worthwhile if you don't plan to reprogram it much as time delay in reprogramming. Also requires special compiler support



(OpenMP?)

- ASIC – can have hard-coded custom hardware for acceleration. Expensive. Found in BitCoin mining?
- DSPs – can be used as accelerators



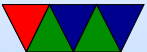
# Why GPUs?

- Old example:
  - 3GHz Pentium 4, 6 GFLOPS, 6GB/sec peak
  - GeForceFX 6800: 53GFLOPS, 34GB/sec peak
- Newer example
  - Raspberry Pi, 700MHz, 0.177 GFLOPS
  - On-board GPU: Video Core IV: 24 GFLOPS



# Key Idea

- using many slimmed down cores
- have single instruction stream operate across many cores (SIMD)
- avoid latency (slow textures, etc) by working on another group when one stalls



# Latency vs Throughput

- CPUs = Low latency, low throughput
- GPUs = high latency, high throughput
- CPUs optimized to try to get lowest latency (caches); with no parallelism have to get memory back as soon as possible
- GPUs optimized for throughput. Best throughput for all better than low-latency for one



# GPU Benefits

- Specialized hardware, concentrating on arithmetic. Transistors for ALUs not cache.
- Fast 32-bit floating point (16-bit?)
- Driven by commodity gaming, so much faster than would be if only HPC people using them.
- Accuracy? 64-bit floating point? 32-bit floating point? 16-bit floating point? Doesn't matter as much if color slightly off for a frame in your video game.
- highly parallel





# GPU Problems

- optimized for 3d-graphics, not always ideal for other things
- Need to port code, usually can't just recompile cpu code.
- Companies secretive.
- serial code
- a lot of control flow
- lot of off-chip memory transfers



# Older / Traditional GPU Pipeline

- In old days, fixed pipeline (lots of triangles).
- Modern chips much more flexible, but the old pipeline can still be implemented in software via the fancier interface.



# Older / Traditional GPU Pipeline

- CPU send list of vertices to GPU.
- Transform (vertex processor) (convert from world space to image space). 3d translation to 2d, calculate lighting. Operate on 4-wide vectors ( $x,y,z,w$  in projected space,  $r,g,b,a$  color space)
- Rasterizer – transform vertexes/vectors into a grid. Fragments. break up to pixels and anti-alias



- Shader (Fragment processor) compute color for each pixel. Use textures if necessary (texture memory, mostly read)
- Write out to framebuffer (mostly write)
- Z-buffer for depth/visibility



# GPGPUs

- Started when the vertex and fragment processors became generically programmable (originally to allow more advanced shading and lighting calculations)
- By having generic use can adapt to different workloads, some having more vertex operations and some more fragment



# Graphics vs Programmable Use

Vertex	Vertex Processing	Data	MIMD processing
Polygon	Polygon Setup	Lists	SIMD Rasterization
Fragment	Per-pixel math	Data	Programmable SIMD
Texture	Data fetch, Blending	Data	Data Fetch
Image	Z-buffer, anti-alias	Data	Predicated Write



# Example for Shader 3.0, came out DirectX9

They are up to Pixel Shader 5.0 now



# Shader 3.0 Programming – Vertex Processor

- 512 static / 65536 dynamic instructions
- Up to 32 temporary registers
- Simple flow control
- Texturing – texture data can be fetched during vertex operations





- Can do a four-wide SIMD MAD (multiply ADD) and a scalar op per cycle:
  - EXP, EXPP, LIT, LOGP (exponential)
  - RCP, RSQ (reciprocal, r-square-root)
  - SIN, COS (trig)



# Shader 3.0 Programming – Fragment Processor

- 65536 static / 65536 dynamic instructions (but can time out if takes too long)
- Supports conditional branches and loops
- fp32 and fp16 internal precision
- Can do 4-wide MAD and 4-wide DP4 (dot product)



# Program

- Typically textures read-only. Some can render to texture, only way GPU can share RAM w/o going through CPU. In general data not written back until entire chunk is done. Fragment processor can read memory as often as it wants, but not write back until done.
- Only handle fixed-point or floating point values
- Analogies:
  - Textures == arrays



- Kernels == inner loops
- Render-to-texture == feedback
- Geometry-rasterization == computation. Usually done as a simple grid (quadrilateral)
- Texture-coordinates = Domain
- Vertex-coordinates = Range



# Flow Control, Branches

- only recently added to GPUs, but at a performance penalty.
- Often a lot like ARM conditional execution



# Terminology (CUDA)

- Thread: chunk of code running on GPU.
- Warp: group of thread running at same time in parallel simultaneously
- Block: group of threads that need to run
- Grid: a group of thread blocks that need to finish before next can be started



# Terminology (cores)

- Confusing. Nvidia would say GTX285 had 240 stream processors; what they mean is 30 cores, 8 SIMD units per core.



# CUDA Programming

- Since 2007
- Use `nvcc` to compile
- `*host*` vs `*device*`
  - host code runs on CPU
  - device code runs on GPU
- Host code compiled by host compiler (`gcc`), device code by custom NVidia compiler





- `__global__` parameters to function – means pass to CUDA compiler
- `cudaMalloc()` to allocate memory and pointers that can be passed in
- call global function like this `add<<<1,1>>>(args)`  
where first inside brackets is number of blocks, second is threads per block
- `cudaFree()` at the end
- Can get block number with `blockIdx.x` and thread index



with `threadIdx.x`

- Can have 65536 blocks and 512 threads (At least in 2010)
- Why threads vs blocks?  
Shared memory, block specific  
`__shared__` to specify
- `__syncthreads()` is a barrier to make sure all threads finish before continuing



# CUDA Programming

- See the NVIDIA “CUDA C Programming Guide”
- Compute Unified Device Architecture
- From CUDA C Programming guide from NVIDIA
- CUDA introduced in 2006
- Heterogeneous programming – there is a host executing a main body of code (a CPU) and it dispatches code to run on a device (a GPU)
- CUDA assumes host and device each have own separate DRAM memory



- CUDA C extends C, define C functions "kernels" that are executed N times in parallel by N CUDA threads



# CUDA Coding

- version compliance – can check version number. New versions support more hardware but sometimes drop old
- nvcc – wrapper around gcc. global code compiled into PTX (parallel thread execution) ISA
- can code in PTX code directly which is sort of like assembly language. Won't give out *actual* assembly language. Why?
- CUDA C has mix of host and device code. Compiles the global stuff to PTX, compiles the <<< ... >>> into



code that can launch the GPU code

- PTX code is JIT compiled into native by the device driver
- You can control JIT with environment variables
- Only subset of C/C++ supported in the device code



# CUDA Hardware

- GPU is array of Streaming Multiprocessors (SMs)
- Program partitioned into blocks of threads that execute independently from each other.
- Manages/Schedules/Executes threads in groups of 32 parallel threads (warps) (weaving terminology) (no relation)
- Threads have own PC, registers, etc, and can execute independently
- When SM given thread block, partitions to warps and



each warp gets scheduled

- One common instruction at a time. If diverge in control flow, each way executed and thread not taking that path just waits.
- Full context stored with each warp; if warp is not ready (waiting for memory) then it may be stopped and another warp that's ready can be run





# CUDA Threads

- kernel defined using `__global__` declaration. When called use `<<<...>>>` to specify number of threads
- each thread that is called is assigned a unique ThreadID  
Use `threadIdx` to find what thread you are and act accordingly

```
__global__ void VecAdd(float *A, float *B)
    int i = threadIdx.x;
    C[i]=A[i]+B[i];
}
```



```

int main(int argc, char **argv) {
    . . . . .
    /* Invoke N threads */
    VecAdd<<<1,N>>>(A, B, C);
}

```

- threadIdx is 3-component vector, can be seen as 1, 2 or 3 dimensional block of threads (thread block)
- Much like our sobel code, can look as 1D (just x), 2D, (thread iD is  $((y * xsize) + x)$  or  $(z * xsize * ysize) + y * xsize + x$ )
- Weird syntax for doing 2 or 3d.



```
--global-- void MatAdd(float A[N][N], flo  
{  
    int i=threadIdx.x;  
    int j=threadIdx.y;  
    C[i][j]=A[i][j]+B[i][j];  
}
```

```
int numBlocks=1;  
dim3 threadsPerBlock(N,N);  
MatAdd<<<numBlocks, threadsPerBlock>>>(A,
```



- Each block made up of the threads. Can have multiple levels of blocks too, can get block number with blockIdx
- Thread blocks operate independently, in any order. That way can be scheduled across arbitrary number of cores (depends how fancy your GPU is)



# CUDA Memory

- Per-thread private local memory
- Shared memory visible to whole block (lifetime of block)
- Global memory
- also constant and texture spaces. Have special rules.  
Texture can do some filtering and stuff
- Global, constant, and texture persistent across kernel launches by same app.



# More Coding

- No explicit initialization, done automatically first time you do something (keep in mind if timing)
- Global Memory: linear or arrays.
  - Arrays are textures
  - Linear arrays are allocated with `cudaMalloc()`, `cudaFree()`
  - To transfer use `cudaMemcpy()`
  - Also can be allocated `cudaMallocPitch()` `cudaMalloc3D()` for alignment reasons



- Access by symbol (?)
- Shared memory, `__shared__`. Faster than Global also `__device__`

Manually break your problem into smaller sizes



# Misc

- Can lock host memory with `cudaHostAlloc()`. Pinned, can't be paged out. Can load store while kernel running if case. Only so much available. Can be marked writecombining. Not cached. So slow for host to read (should only write) but speeds up PCI transaction.





# Async Concurrent Execution

- Instead of serial/parallel/serial/parallel model
- Want to have CUDA running and host at same time, or with mem transfers at same time
  - Concurrent host/device: calls are async and return to host before device done
  - Concurrent kernel execution: newer devices can run multiple kernels at once. Problem if use lots of memory
  - Overlap of Data Transfer and Kernel execution
  - Streams: sequence of commands that execute in order,



but can be interleaved with other streams  
complicated way to set them up. Synchronization and  
callbacks



# Events

- Can create performance events to monitor timing
- PAPI can read out performance counters on some boards
- Often it's for a full synchronous stream, can't get values mid-operation
- NVML can measure power and temp on some boards?



# Multi-device system

- Can switch between active device
- More advanced systems can access each others device memory



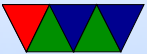
# Other features

- Unified virtual address space (64 bit machines)
- Interprocess communication
- Error checking



# Texture Memory

- Complex



# 3D Interop

- Can make results go to an OpenGL or Direct3D buffer
- Can then use CUDA results in your graphics program



# Code Example

```
#include <stdio.h>
```

```
#define N 10
```

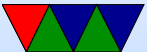
```
__global__ void add (int *a, int *b, int *c)  
    int tid=blockIdx.x;  
  
    if (tid < N) {  
        c[tid]=a[tid]+b[tid];  
    }
```





```
    }  
}
```

```
int main(int argc, char **argv) {  
  
    int a[N], b[N], c[N];  
    int *dev_a, *dev_b, *dev_c;  
    int i;  
  
    /* Allocate memory on GPU */
```



```
cudaMalloc (( void **) &dev_a , N * sizeof ( int ) ) ;
cudaMalloc (( void **) &dev_b , N * sizeof ( int ) ) ;
cudaMalloc (( void **) &dev_c , N * sizeof ( int ) ) ;

/* Fill the host arrays with values
for ( i = 0 ; i < N ; i ++ ) {
    a [ i ] = - i ;
    b [ i ] = i * i ;
}

cudaMemcpy ( dev_a , a , N * sizeof ( int ) , cu
```



```
cudaMemcpy( dev_b , b , N* sizeof( int ) , cu
add<<<N,1>>>( dev_a , dev_b , dev_c );
cudaMemcpy( c , dev_c , N* sizeof( int ) , cu

/* results */
for( i=0; i<N; i++) {
    printf( " %d+ %d= %d \n" , a[ i ] , b
}
```



```
    cudaFree ( dev_a );  
    cudaFree ( dev_b );  
    cudaFree ( dev_c );  
  
    return 0;  
}
```

