

# **ECE 571 – Advanced Microprocessor-Based Design Lecture 3**

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# Announcements

- HW#1 will be posted
- I will be handing out account username/passwords for the homework.
- Accounts: Log in to the haswell “haswell-ep” machine for homework. Make sure you connect to port 2131. ece571-1 names are a bit impersonal.  
Use `passwd` to change your password.  
You can use `chfn` to change your name as it appears in `w` if you want.



Please use the accounts wisely

- Don't forget no class Monday, labor day



# Measuring Performance



# Using time

- For example

```
$ time xhpl
...
real    0m9.484s
user    0m29.150s
sys     0m7.440s
```

Real Time = Wall clock

User Time = Time used by program alone

Sys Time = Time used by OS

- When could Real be greater than User?



Other users/jobs on system.

When could User be greater than Real?

Multiple threads.

- Run multiple times and notice time changes



# What if Time isn't Enough?



# What are Hardware Performance Counters?

- Registers on CPU that measure low-level system performance
- Available on most modern CPUs; increasingly found on GPUs, network devices, etc.
- Low overhead to read





# Hardware Implementation of Counters

- Not much documentation available
- Jim Callister/Intel: “Confessions of a Performance Monitor Hardware Designer” 2005, Workshop on Hardware Performance Monitor Design
  - Transistors free, wires not. Also design time, validation, documentation, time to market. PMU has tentacles “everywhere” bringing data back to center.
  - Architect too much, lower performance, events don’t



map well to hardware. Architect too little.. software design harder.

- Which events are important? Are cache misses important if don't hurt performance? (no stalls)
- Mapping events to signal difficult. On critical path. Not enough wires. Combining signals hard if distance between wires.
- Use logging. May miss events in “shadow” of another event being logged. Use random behavior?



# Learning About the Counters

- Number of counters varies from machine to machine
- Available events different for every vendor and every generation
- Available documentation not very complete (Intel Vol3b, AMD BKDG, ARM ARM/TRM)



# Low-level interface

- on x86: MSRs
- ARM: CP15 system control register



# Overflow

- overflows after hitting a threshold (often when wrapping, most counters are between 32 and 44 bits wide)
- One use is to keep track of counters that may wrap multiple times between reads
- If want to overflow earlier, init to a high value. So `0xc0000000` to overflow at 1 billion



# Accuracy, Determinism vs Overcount

- Determinism – same count every time you run
- Overcount – an event counts more than the expected amount



# SW Sources of Non-Determinism

- Accessing changing values, such as time
- Pointer-value dependencies



# Linux interface

- Abstract away.
- `perf_event_open()`. See the manpage.
- Very complicated system call.
- Most people use `perf` or `PAPI` rather than calling it directly.





# perf tool

A a tutorial on perf can be found here:

<https://perf.wiki.kernel.org/index.php/Tutorial>

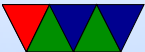


# perf list

Lists available events

List of pre-defined events (to be used in `-e`):

<code>cpu-cycles OR cycles</code>	<code>[Hardware event]</code>
<code>instructions</code>	<code>[Hardware event]</code>
<code>cache-references</code>	<code>[Hardware event]</code>
<code>cache-misses</code>	<code>[Hardware event]</code>
<code>branch-instructions OR branches</code>	<code>[Hardware event]</code>
<code>branch-misses</code>	<code>[Hardware event]</code>
<code>bus-cycles</code>	<code>[Hardware event]</code>
<code>cpu-clock</code>	<code>[Software event]</code>
<code>task-clock</code>	<code>[Software event]</code>
<code>page-faults OR faults</code>	<code>[Software event]</code>
<code>minor-faults</code>	<code>[Software event]</code>
<code>major-faults</code>	<code>[Software event]</code>
<code>context-switches OR cs</code>	<code>[Software event]</code>



# perf stat – Aggregate results

```
vince@arm:~/class/ece571$ perf stat ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
```

```
Performance counter stats for './matrix_multiply':
```

```
11585.144036 task-clock # 0.999 CPUs utilized
      19 context-switches # 0.000 M/sec
      0 CPU-migrations # 0.000 M/sec
    1,633 page-faults # 0.000 M/sec
10,343,746,076 cycles # 0.893 GHz
    5,031,717 stalled-cycles-frontend # 0.05% frontend cycles idle
    9,521,135,479 stalled-cycles-backend # 92.05% backend cycles idle
    1,176,286,814 instructions # 0.11 insns per cycle
                                # 8.09 stalled cycles per insn
    137,835,961 branches # 11.898 M/sec
      831,736 branch-misses # 0.60% of all branches

11.591796875 seconds time elapsed
```



# perf stat – Specifying Events

```
vince@arm:~/class/ece571$ perf stat -e instructions,cycles ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
```

```
Performance counter stats for './matrix_multiply':
```

```
1,174,788,622 instructions          #    0.14  insns per cycle
8,346,588,065 cycles                #    0.000 GHz
```

```
12.394775391 seconds time elapsed
```



# perf stat – Specifying Masks

:u is user, :k kernel

ARM Cortex A9 cannot specify this distinction (results shown here are x86)

```
vince@arm:~/class/ece571$ perf stat -e instructions,instructions:u ./matri
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix_multiply':

   950,526,051 instructions          #    0.00  insns per cycle
   945,661,967 instructions:u       #    0.00  insns per cycle

1.052072277 seconds time elapsed
```



# libpfm4 – Finding All Event Names

```
./showevtinfo
Supported PMU models:
    [51, perf, "perf_events generic PMU"]
    [65, arm_ac8, "ARM Cortex A8"]
    [66, arm_ac9, "ARM Cortex A9"]
    [75, arm_ac15, "ARM Cortex A15"]
Detected PMU models:
    [51, perf, "perf_events generic PMU", 80 events, 1 max encoding, 0 counters, OS g
    [66, arm_ac9, "ARM Cortex A9", 57 events, 1 max encoding, 2 counters, core PMU]
Total events: 254 available, 137 supported
...
#-----
IDX      : 138412068
PMU name : arm_ac9 (ARM Cortex A9)
Name     : NEON_EXECUTED_INST
Equiv    : None
Flags    : None
Desc     : NEON instructions going through register renaming stage (approximate)
Code     : 0x74
#-----
....
```



# libpfm4 – Finding Raw Event Values

```
./check_events NEON_EXECUTED_INST
Supported PMU models:
[51, perf, "perf_events generic PMU"]
[65, arm_ac8, "ARM Cortex A8"]
[66, arm_ac9, "ARM Cortex A9"]
[75, arm_ac15, "ARM Cortex A15"]
Detected PMU models:
[51, perf, "perf_events generic PMU"]
[66, arm_ac9, "ARM Cortex A9"]
Total events: 254 available, 137 supported
Requested Event: NEON_EXECUTED_INST
Actual      Event: arm_ac9::NEON_EXECUTED_INST
PMU         : ARM Cortex A9
IDX         : 138412068
Codes      : 0x74
```



# perf – Using Raw Event Values

```
vince@arm:~/class/ece571$ perf stat -e r74 ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
```

```
Performance counter stats for './matrix_multiply':
```

```
1 r74
```

```
11.303955078 seconds time elapsed
```





# perf stat – multiplexing

```
perf stat -e instructions,instructions,branches,cycles,cycles ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094

Performance counter stats for './matrix_multiply':

   1,178,121,057 instructions #    0.12  insns per cycle [40.23%]
   1,180,460,368 instructions #    0.12  insns per cycle [60.25%]
     138,550,072 branches                               [80.09%]
   9,999,614,616 cycles #    0.000 GHz                [79.85%]
   9,926,949,659 cycles #    0.000 GHz                [20.17%]

11.214630127 seconds time elapsed
```

Note same event not same results, approximate because an estimate. Percentage shown is percentage event was active during run.



# perf stat – all cores

```
vince@arm:~/class/ece571$ sudo perf stat -a ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
```

```
Performance counter stats for './matrix_multiply':
```

```
24089.660644 task-clock                #    2.001 CPUs utilized          [100.00%]
      105 context-switches             #    0.000 M/sec                   [100.00%]
      1,641 page-faults                 #    0.000 M/sec                   [100.00%]
9,218,451,619 cycles                    #    0.383 GHz                     [100.00%]
      9,707,195 stalled-cycles-frontend #    0.11% frontend cycles idle   [100.00%]
8,393,095,067 stalled-cycles-backend   #   91.05% backend cycles idle    [100.00%]
1,193,164,945 instructions              #    0.13 insns per cycle         [100.00%]
                                           #    7.03 stalled cycles per insn [100.00%]
139,913,572 branches                   #    5.808 M/sec                   [100.00%]
      1,221,237 branch-misses          #    0.87% of all branches        [100.00%]

12.040527344 seconds time elapsed
```

Run on *all* cores of system even if your process not running there. `-a` option. Need root permissions. (Why? Security)



# perf record – sampling

```
vince@arm:~/class/ece571$ time ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094

real0m10.747s
user0m10.688s
sys0m0.055s
vince@arm:~/class/ece571$ time perf record ./matrix_multiply
Matrix multiply sum: s=27665734022509.746094
[ perf record: Woken up 2 times to write data ]
[ perf record: Captured and wrote 0.454 MB perf.data (~19853 samples) ]

real0m12.009s
user0m11.797s
sys0m0.203s
```

perf record creates perf.data, use -o to specify output



# perf report – summary of recorded data

```
99.62% matrix_multiply matrix_multiply      [.] naive_matrix_multiply
0.38%  matrix_multiply [kernel.kallsyms].head.text [k] 0xc0046a54
0.00%  matrix_multiply ld-2.13.so          [.] _dl_relocate_object
0.00%  matrix_multiply [kernel.kallsyms]          [k] __do_softirq
```

Our benchmark is simple (only one function) so the profiled results are not that exciting.

The [k] indicates that profile happened while the kernel was running.



# Similar ways to get Similar Results

- Valgrind/Callgrind

*valgrind - -tool=callgrind BENCHMARK*

then run *callgrind\_annotate*

Note Valgrind is probably around 50 times slower

- Use gprof

Compile your code with *-pg*

Run *gprof BENCHMARK*



# perf annotate – show hotspots in assembly

```
0.00 :          845a:      vldr    d7, [pc, #124] ; 84d8 <naive_matrix_m
30.97 :          845e:      adds   r1, r4, r3
1.43 :          8460:      add.w  r3, r3, #4096 ; 0x1000
1.17 :          8464:      adds   r2, #8
1.36 :          8466:      cmp.w  r3, #2097152 ; 0x200000
2.97 :          846a:      vldr   d5, [r2]
2.62 :          846e:      vldr   d6, [r1]
2.78 :          8472:      mov    r9, r2
2.42 :          8474:      vmla.f64      d7, d5, d6
53.81 :          8478:      bne.n  845e <naive_matrix_multiply+0x72>
0.01 :          847a:      adds   r5, #1
```

The annotated results show a branch and an add instruction accounting for 83% of profiles. Likely this is due to skid and the key instruction is the previous `vmla.f64` floating point multiply instruction. The processor just isn't able to stop at the exact instruction when the interrupt comes in.



# Skid

- Beware of “skid” in sampled results
- This is what happens when a complex processor cannot stop immediately, so the reported instruction might be off by a few instructions.
- Some processors do not have this problem, other Intel processors have special events that can compensate for this.

