

ECE 571 – Advanced Microprocessor-Based Design Lecture 15

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Announcements

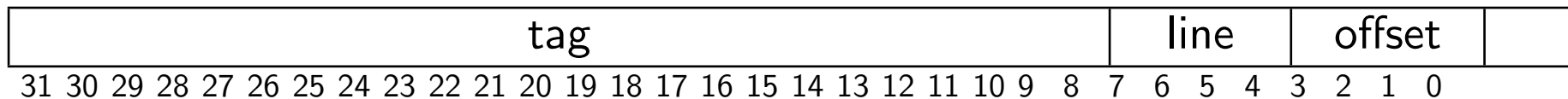
- HW#5 will be posted, caches



Cache Example Two

512 Byte cache, 2-Way Set Associative, with 16 byte lines, LRU replacement.

24-bit tag, 16 lines (4 bits), 4-bit offset.



Cache Example 2



Cache Example – Instruction 1

```
ldrb r1, 0x00000000
```

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	0	0	0000 00	0			
1	0				0			
2	0				0			
3	0				0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Miss, Cold



Cache Example – Instruction 2

`str r1, 0x00000001`

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	1	0	0000 00	0			
1	0				0			
2	0				0			
3	0				0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Hit



Cache Example – Instruction 3

strb r1, 0x00000105

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	1	1	0000 00	1	1	0	0000 01
1	0				0			
2	0				0			
3	0				0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Miss, Cold



Cache Example – Instruction 4

```
ldr r1, 0x00000206
```

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	0	0	0000 02	1	1	1	0000 01
1	0				0			
2	0				0			
3	0				0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Miss, Cold



Cache Example – Instruction 5

ldb r1, 0x00000000

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	0	1	0000 02	1	0	0	0000 00
1	0				0			
2	0				0			
3	0				0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Miss, Conflict



Cache Example – Instruction 6

ldb r1, 0x00000030

	Way 0				Way 1			
line	V	D	LRU	Tag	V	D	LRU	Tag
0	1	0	1	0000 02	1	0	0	0000 00
1	0				0			
2	0				0			
3	1	0	0	0000 00	0			
4	0				0			
5	0				0			
...								
b	0				0			
c	0				0			
d	0				0			
e	0				0			
f	0				0			

Miss, Cold



CMP Issues



Cache Coherency

- Protocols such as MESI (Modified, Exclusive, Shared, Invalid)
- Snoopy vs Directory

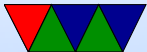
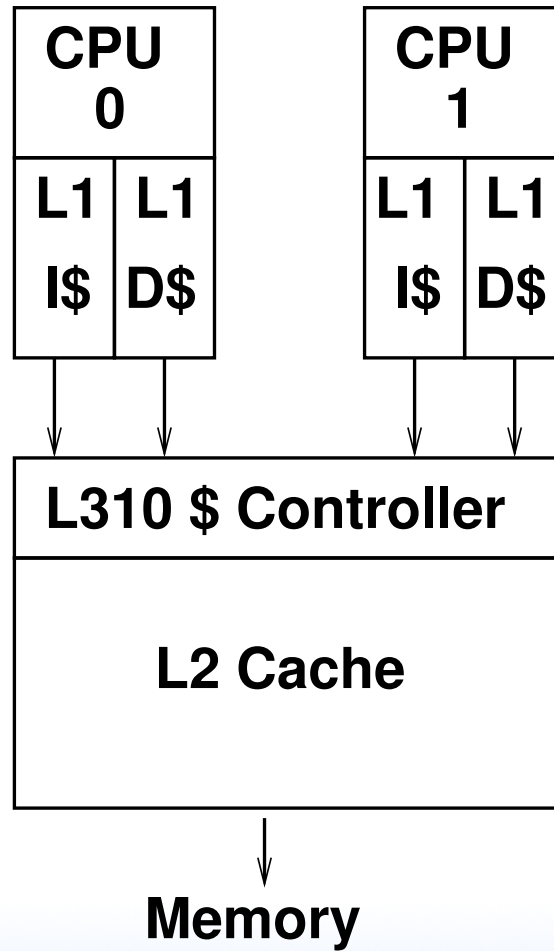


MESI Example

- Only gets complicated when you start writing.
- Invalid – starts out here
- Shared – only has been read, same as memory, can be in multiple caches
- Exclusive – a core is requesting to read, so it gets exclusive access, invalidates all other copies
- Modified – dirty, has been written to. Write back and then can change to S



Cortex A9 Cache Layout



Cortex A9 Cache Layout

- OMAP4430 processor
- 32kB 4-way associative, separate L1-I and L1-D
 - pseudo-round-robin or pseudo random replacement
 - 8-word line size (32B)
 - critical-word first filling
 - instruction: VIPT, data: PIPT
- Optional L2 cache controller
 - Pandaboard has L310 L2 cache controller, 1MB 16-way



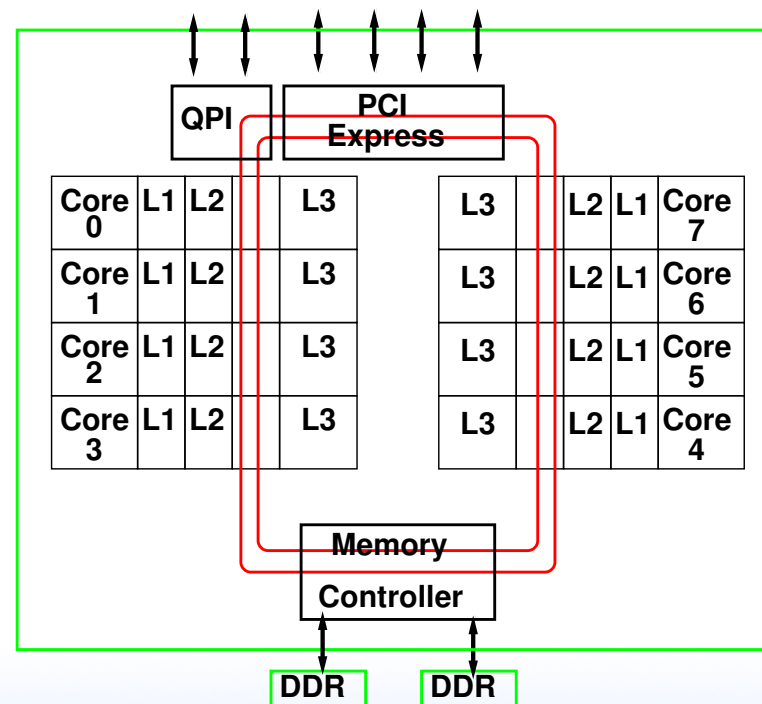
Optional prefetcher

- data cache reads/writes non-blocking, 4 outstanding misses
write buffer: 4 64-bit, allowing write combining

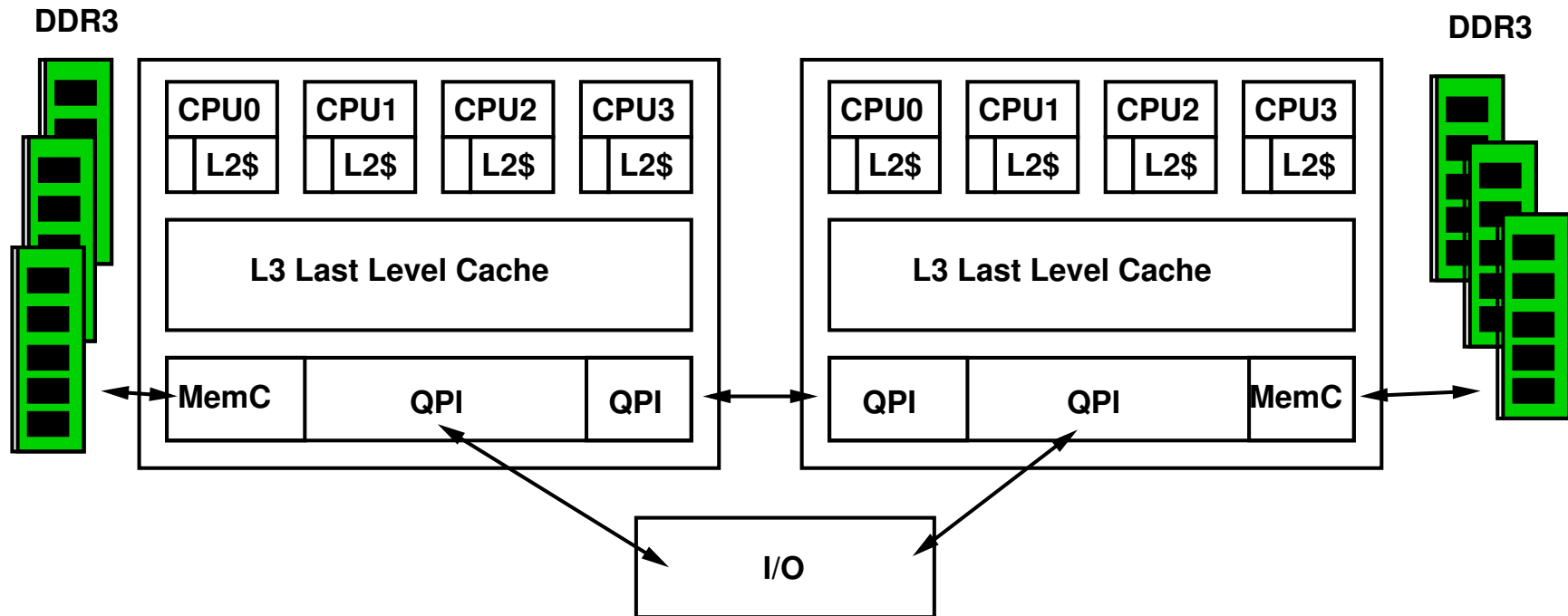


Haswell-EP Cache Layout

Note: see “Cache Coherence Protocol and Memory Performance of the Intel Haswell-EP Architecture” by Molka, Hackenberg, Schöne, Nagel.



SandyBridge Cache Layout



SandyBridge Cache Layout

- per core 32kB L1 I/D – 4 clocks
64B/line, 8-Way
(shared if hyper-threaded)
writeback
- mOp cache? 1.5K instructions, 8-way, 6Mop/line
Loop stream detector, can execute w/o accessing icache
- per core 256kB L2 unified – 12 clocks



64B/line, 8-way
writeback. non-inclusive

- shared L3 1MB-20MB – 26-31 clocks
64B/line. 12-way (varies)
writeback, inclusive
- various hw prefetchers operating

