# ECE 574 – Cluster Computing Lecture 9

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29 September 2015

#### Announcements

 Homework #3 was posted.
 Sorry for delay, Mainestreet down Friday night and most of Saturday



## Go Over HW#3

• Convolution, good example found in Wikipedia:

https://en.wikipedia.org/wiki/Kernel\_%28image\_processing%29

- libjpeg and layout. Makes an array of size x\*y\*3 bytes.
   The three bytes (24-bits) per pixel are r,g,b
- C arrays are complex. 3D array pixels[color][xsize][ysize] can also (and sometimes can only) be accesses as equivalent 1D array:

pixels[d][x][u]=pixels+((y\*xsize\*3)+(x\*3)+d);

sobelx, sobely, sobel



- There was a minor bug in the provided sample output, was forgetting to saturate in the final step.
- Edge handling:

extend – the edge pixel is extended out

wrap – wraps around to other side of image

crop – output is smaller by 1 pixel border (alternately, just fill it with 0s)

any of these is fine, my example crops and fills with 0s

PAPI – last level cache misses.
 How many should there be?
 320x320x3bytes = 300kb.



- in, outx, outy, out total = 300kB\*4 = 1.2MB
- Most likely fits in LLC (8MB on this machine, see /proc/cpuinfo).
- In theory cache misses would be 1.2MB/line size (64 bytes)
- is 19200 misses in theory. Why might it be lower?
- Something cool just make some optimization, does not need to be successful. Compare timing and L3 Total Cache Misses.



#### OpenMP

A good writeup is here: https://computing.llnl.gov/tutorials/openMP/



## OpenMP

- Shared memory multi-processing interface
- $\bullet$  C, C++ and FORTRAN
- Industry standard made by lots of companies
- gcc support fairly recently donated, CLANG support even more recent
- OpenMP 1.0 came out in 1997 (FORTRAN) or 1998 (C), now version 4.0 (2013)



gcc added support in 4.2 (OpenMP 2.5)
4.4 (OpenMP 3.0), 4.7 (OpenMP 3.1), 4.9 (OpenMP 4.0), 5.0 (Offloading)



## OpenMP

- Based on threads
- Master thread with Fork/Join methodology Master thread forks off threads which compute in parallel.
  - When the threads complete, then join back to the master thread.
  - You can repeat many cycles of this.
- Include diagram



- Indicate parallelism with #pragma omp
- Can possibly have nested threads (implementation dependent).
- Can possibly have dynamic num of threads (implementation dependent)
- Relaxed consistency, threads can cache local variables, so if you need memory to be consistent might need to flush it.



### **OpenMP Interface**

- Compiler Directives
- Runtime Library Routines
- Environment Variables



## **Compiler Support**

- On gcc, pass -fopenmp
- C: #pragma omp
- FORTRAN: C\$OMP or !\$OMP



### **Compiler Directives**

- Spawning a parallel region
- Dividing blocks of code among threads
- Distributing loop iterations between threads
- Serializing sections of code
- Synchronization of work among threads



### Library routines

- Need to #include <omp.h>
- Getting and setting the number of threads
- Getting a thread's ID
- Getting and setting threads features
- Checking in in parallel region
- Checking nested parallelism



- Locking
- Wall clock measurements



#### **Environment Variables**

- Setting number of threads
- Configuring loop iteration division
- Processor bindings
- Nested Parallelism settings
- Dynamic thread settings
- Stack size



• Wait policy



#### Simple Example

```
#include <stdio.h>
#include <stdlib.h>
#include <omp.h>
int main (int argc, char **argv) {
        int nthreads,tid;
/* Fork a parallel region, each thread having private copy of tid */
#pragma omp parallel private(tid)
        {
        tid=omp_get_thread_num();
        printf("\tInside of thread %d\n",tid);
        if (tid==0) {
                nthreads=omp_get_num_threads();
                printf("This is the master thread, there are %d threads\n",
                        nthreads);
        }
```



}

}

/\* End of block, waits and joins automatically \*/

return 0;



#### Notes on the example

- PARALLEL directive creates a set of threads and leaves the original thread the master, with tid 0.
- All threads will execute the code in parallel region
- There's an implied barrier/join at end of parallel region. Only the master continues after it.
- If any thread terminates in a parallel region, then all threads will also terminate.
- You can't goto into a parallel region.



#### parallel clause

structured\_block

- if you can do a check if (i==0)
   If true parallel threads are created, otherwise serially
- private variables in the list are private to each thread



- shared variables in the list shared across threads
- default shared or none (more on C), default scope.
   none means you have to explicitly share or private each var
- firstprivate a variable inside a parallel section ends up with the value it had before the parallel section
- lastprivate a variable outside the parallel section gets the value from the last loop iteration
- copyin copies in the value from the master thread into



each thread (how is this different from firstprivate?)

- reduction vector dot product. The work is split up into equal chunks, then the operator provided is used to ? and then they are all combined for final result.
- num\_threads number of threads to use



#### How many threads?

- Evaluation of the IF clause
- Setting of the NUM\_THREADS clause
- Use of the omp\_set\_num\_threads() library function
- Setting of the OMP\_NUM\_THREADS environment variable
- Implementation default usually the number of CPUs on a node, though it could be dynamic (see next bullet).



• Threads are numbered from 0 (master thread) to N-1



#### **Work-sharing Constructs**

- Must be inside of a parallel directive
- $\bullet$  do/for
- sections
- single



# Do/For

```
#pragma omp for [clause ...] newline
schedule (type [,chunk])
ordered
private (list)
firstprivate (list)
lastprivate (list)
shared (list)
reduction (operator: list)
collapse (n)
nowait
```

for\_loop

• schedule

static – divided into size chunk, statically assigned to threads dynamic – divided into chunks, dynamically



assigned threads as they finish guided – like dynamic but shrinking blocksize runtime – from OMP\_SCHEDULE environment variable auto – compiler picks for you

- nowait threads do not wait at end of loop
- ordered loops must execute in order they would in serial code
- collapse nested loops can be collapsed?



#### **For Example**

```
#include <stdio.h>
#include <stdlib.h>
#include <omp.h>
static char *memory;
int main (int argc, char **argv) {
    int num_threads=1;
    int mem_size=256*1024*1024; /* 256 MB */
    int i,tid,nthreads;
    /* Set number of threads from the command line */
    if (argc>1) {
        num_threads=atoi(argv[1]);
    }
    /* allocate memory */
    memory=malloc(mem_size);
    if (memory==NULL) perror("allocating memory");
```



```
#pragma omp parallel shared(mem_size,memory) private(i,tid)
{
    tid=omp_get_thread_num();
    if (tid==0) {
        nthreads=omp_get_num_threads();
        printf("Initializing %d MB of memory using %d threads\n",
            mem_size/(1024*1024),nthreads);
    }
    #pragma omp for schedule(static) nowait
    for (i=0; i < mem_size; i++)
        memory[i]=0xa5;
    }
    printf("Master thread exiting\n");
</pre>
```



}